Results 000000000000 Purification in Normal Form Games

Purification in Dynamic Games

Lecture 14: Limited Memories & Purification

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Overview

Last lecture: Reputation and social learning.

- Each short-run player observes a bounded number of the long-run player's action, and at least one short-run player's action.
- Proof: A mixed-strategy equilibrium in which learning is slow.

This lecture:

- Reputation effects with limited memory w/o social learning.
- Critique of mixed-strategy equilibria in extensive-form games with limited memories.

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Primitives

- Long-lived player 1 vs short-lived player 2s.
- Player 1's action $x \in [0, 1]$. Player 2's action $y \in [0, 1]$.
- Player 1's stage-game payoff $u_1(x, y)$ satisfies:
 - 1. $u_1(x, y)$ is strictly decreasing in x,
 - 2. $u_1(x, y)$ is strictly increasing in y,
 - 3. $u_1(x, y)$ has strictly decreasing differences in (x, y).
- P2's stage-game payoff $u_2(x, y)$ has strictly increasing differences, P2 has a unique best reply to any of P1's mixed actions.
 - P2's best reply y* : Δ[0, 1] → [0, 1] is continuous and strictly increasing under FOSD.
 - Let us normalize $y^*(0) = 0$.
- P1's dominant action is 0. Unique stage-game NE is (0,0).

Results 000000000000

Long-Run Player's Type

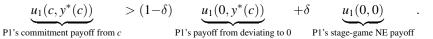
Player 1 has two types:

- With prob $\mu^* \in (0, 1)$, P1 is a commitment type who plays c > 0.
- With prob $1 \mu^*$, P1 is rational and maximizes discounted payoff.

We assume that commitment to *c* is valuable:

$$u_1(c, y^*(c)) > u_1(0, 0).$$

We also assume that δ is large enough such that:



Monitoring Structure

 $P2_t$ only observes P1's action in the last K periods.

•
$$(a_{\max\{0,t-K\}},...,a_{t-1}).$$

Important things to emphasize:

- Player 2 cannot observe the actions of previous short-run players.
- Player 2 cannot observe calendar time.

No social learning.

Player 2's prior assigns prob $(1 - \delta)\delta^t$ to calendar time being *t*, and updates her belief according to Bayes rule after observing $(a_{\max\{0,t-K\}}, ..., a_{t-1})$.

Strategies & Stationary Perfect Bayesian equilibrium

The set of P2's histories $\mathcal{H} \equiv \bigcup_{k=0}^{K} [0, 1]^k$.

• Sequences of player 1's actions with length no more than *K*.

P2's strategy $\sigma_2 : \mathcal{H} \to [0, 1]$.

Focus on *stationary equilibria* s.t. P1's strategy depends only on \mathcal{H} .

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Sufficient Statistics Result

Partition the set of length *K* histories into K + 1 subsets:

- $h \in \mathcal{H}^K$ if the last *K* actions were *c*. Call them clean histories.
- For every h ∈ H\H^K, let I(h) ∈ {0, 1, ..., K − 1} be the number of periods since the most recent action that is not c.

Intuitively, if I(h) is larger, then h is closer to a clean history.

Let \mathcal{H}^k be the set of histories with I(h) = k, for $k \in \{0, 1, ..., K - 1\}$.

Proposition: Sufficient Statistics Result

In any stationary equilibrium, (i) players' actions after period K are measurable with respect to the above partition, and (ii) player 1 will only play 0 and c with positive probability.

This depends on the submodularity of P1's payoff.

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Proof Sketch: Sufficient Statistics Result

Consider two histories h and h' s.t.

• P1's actions were the same in the last K - 1 periods,

P1's action K periods ago was not c under h and h'.

Lemma

P2's actions at h and h' are the same. P1's continuation values at h and h' are the same.

Suppose by way of contradiction that $a_2(h) > a_2(h') \ge 0$.

- Since h and h' are not clean, P1 is known to be rational.
- $a_2(h)$ depends only on $a_1(h)$, $a_2(h')$ depends only on $a_1(h')$.

Hence, there exist a_1^* and a_1' with $a_1^* > a_1'$ s.t.

- P1 plays a_1^* with positive prob at *h*.
- P1 plays a'_1 with positive prob at h'.

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Proof Sketch: Sufficient Statistics Result

Suppose by way of contradiction that $a_2(h) > a_2(h') \ge 0$.

P1 plays a_1^* at h, and plays $a_1'(< a_1^*)$ at h'.

P1's incentive constraint at *h*:

$$(1-\delta)u_1(a_1^*,a_2(h)) + \delta V(h,a_1^*) \ge (1-\delta)u_1(a_1',a_2(h)) + \delta V(h,a_1').$$

P1's incentive constraint at h':

~

$$(1-\delta)u_1(a'_1,a_2(h')) + \delta V(h',a'_1) \ge (1-\delta)u_1(a^*_1,a_2(h')) + \delta V(h',a^*_1).$$

Since *h* and *h'* differ only in action *K* periods ago, we have $V(h, a_1^*) = V(h', a_1^*)$ and $V(h, a_1') = V(h', a_1')$. Hence,

$$u_1(a_1^*, a_2(h)) - u_1(a_1', a_2(h)) \ge \frac{\delta}{1-\delta}(V(h, a_1') - V(h, a_1^*))$$

$$= \frac{\delta}{1-\delta}(V(h',a_1') - V(h',a_1^*)) \ge u_1(a_1^*,a_2(h')) - u_1(a_1',a_2(h'))$$

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Proof Sketch: Sufficient Statistics Result

Therefore, $a_2(h) > a_2(h'), a_1^* > a_1'$, but

 $u_1(a_1^*,a_2(h)) - u_1(a_1',a_2(h)) \ge u_1(a_1^*,a_2(h')) - u_1(a_1',a_2(h')).$

This leads to a contradiction since u_1 has strictly decreasing differences.

Hence, P2's actions at h and h' must be the same.

P1's continuation values at h and h' must be the same as well.

• Why? P2's actions are the same, and P1's continuation value in the next period does not depend on whether the current history is *h* or *h*'.

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Proof Sketch: Sufficient Statistics Result

We have just shown that at any two histories h and h' s.t.

• P1's actions were the same in the last K - 1 periods,

P1's action K periods ago was not c under h and h'.

then

- P2's actions are the same at *h* and *h'*.
- P1's continuation values at the same at *h* and *h'*.

What about h and h' s.t.

• P1's actions were the same in the last K - 2 periods,

P1's action K - 1 periods ago was not c under h and h'?

What can we say about $V(h, a_1)$ and $V(h', a_1)$?

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Proof Sketch: Sufficient Statistics Result

Therefore, for any $k \in \{1, 2..., K\}$, if *h* and *h'* are such that:

• P1's actions were the same in the last k - 1 periods,

P1's action k periods ago was not c under h and h'.

At h and h', P1's continuation values and P2's actions are the same.

Recall the definition of I(h).

If *I*(*h*) = *I*(*h*'), then either *h* = *h*' or there exists *k* ∈ {1, 2, ..., *K*} such that P1's actions were the same in the last *k* − 1 periods and their actions *k* periods ago were not *c*.

Hence, P2's action and P1's continuation value depend only on I(h).

- This implies that P1 plays either 0 or *c*.
- Hence, P1's mixed actions can be ranked according to FOSD.
- If a₂(h) = a₂(h') while h and h' are not clean, then P1's action at h and h' must be the same.

Characterize Stationary PBE

- μ_K : Prob P2's belief assigns to the commitment type at clean histories.
- β_k : Prob that rational-type P1 plays *c* at \mathcal{H}^k , for every $k \in \{0, 1, ..., K\}$.
- y_k : P2's action at \mathcal{H}^k , for every $k \in \{0, 1, ..., K\}$.

Theorem 1

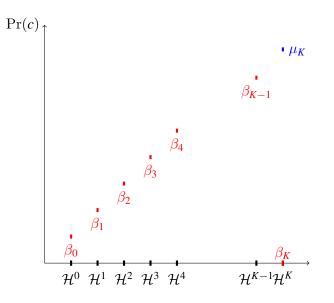
Under any $\delta > \overline{\delta}$, memory length $K \in \mathbb{N}$, and prior belief $\mu^* > 0$. Every stationary PBE takes the following form:

- The rational type P1 plays 0 for sure at clean histories.
- Prob of c increases with the index $0 < \beta_0 < \beta_1 < ... < \beta_{K-1} < \mu_K$.
- Trust increases with the index $y_0 < y_1 < ... < y_{K-1} < y_K$,

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Long-Run Player's Equilibrium Actions (in red)



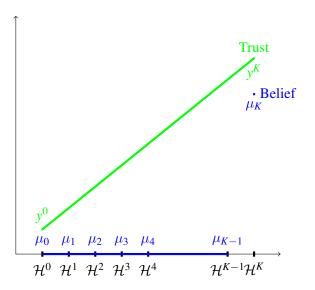
Liu and Skrzypacz Results

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The Short-Run Players' Actions and Beliefs



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Characterize Stationary PBE

Takeaway from their characterization result:

- 1. P2's trust increases with the index I(h).
- 2. P1 betrays for sure at clean histories.
- 3. At non-clean histories, P1's prob of playing c increases with I(h).

Intuition: $u_1(a, b)$ has strictly decreasing differences.

• P1 has stronger incentive to betray when he is trusted more.

Caveat: Supermodularity/submodularity in the stage game usually do not imply much in the repeated game.

• Why? My action today affects your observation tomorrow.

e.g., in the proof of the sufficient statistics result, we only use submodularity at specific pairs of histories.

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Lower Bound on P1's Equilibrium Payoff

Theorem 2

For any $\varepsilon > 0$ and $\mu^* \in (0, 1)$, there exists $K(\varepsilon, \mu^*) \in \mathbb{N}$ such that for any

 $K > K(\varepsilon, \mu^*)$, P1's payoff at any history of any stationary PBE is at least

$$(1-\delta^K)g_1(c,0)+\delta^Kg_1(c,y^*(c))-\varepsilon.$$

When K is large enough,

• as long as P1 is sufficiently patient, he can secure his commitment payoff $g_1(c, y^*(c))$ in *all* stationary PBEs.

This stands in contrast to Fudenberg and Levine (1989,1992):

• No meaningful payoff lower bound that applies to all histories.

Proof Sketch

The proof is different from the one in Fudenberg and Levine:

- Fudenberg and Levine's technique requires P2 observing everything her predecessors observed.
- This is not the case when memories are bounded.

Key step of the proof:

Lemma

For any $\eta \in (0, 1)$, $\exists \overline{K}(\eta)$ such that when $K > \overline{K}(\eta)$, $\mu(K) > 1 - \eta$.

Intuition: Suppose $\mu(K) \leq 1 - \eta$. Since $\beta(k) < \mu(K)$ for every $k \leq K - 1$, the rational type reaches the clean history with prob no more than $(1 - \eta)^K$.

• If K is large, then
$$(1 - \eta)^K$$
 is small,

which implies that P2's belief at clean histories assigns a high prob to the commitment type. This contradicts $\mu(K) \leq 1 - \eta$.

Dynamic Games with Limited Memories

Some of the results rely on constructing mixed-strategy equilibria:

• Liu (2011), Liu and Skrzypacz (2014), Pei (2023).

Mixed-strategy equilibria are also used to show folk theorems in repeated games with private monitoring.

 Most notably, the literature on belief-free equilibria, e.g., Ely and Välimäki (2002), Ely, Hörner and Olzewski (2005).

How robust are these mixed-strategy equilibria?

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Background: Mixed Strategies & Harsanyi Purification

Battle of sexes game:

-	L	R
Т	2, 1	0,0
В	0,0	1,2

There exists one mixed-strategy Nash equilibrium:

• P1 chooses $\frac{2}{3}T + \frac{1}{3}B$. P2 chooses $\frac{1}{3}L + \frac{2}{3}R$.

How to pin down players' mixing probabilities?

- P1's mixing prob is pinned down by P2's indifference condition.
- P2's mixing prob is pinned down by P1's indifference condition.

This logic sounds weird.

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Background: Mixed Strategies & Harsanyi Purification

Harsanyi: What if players privately observe private payoff shocks?

-	L	R
Т	2, 1	0,0
В	0, 0	1,2

Suppose P1's payoff from T equals his payoff in the matrix plus ξ .

• ξ is drawn from a continuous distribution $U[0, \varepsilon]$.

Only P1 observes the realization of ξ .

Suppose P2's payoff from L equals his payoff in the matrix plus η .

η is drawn from a continuous distribution U[0, ε].
 Only P2 observes the realization of η.

P1 chooses T when ξ is high and chooses B when ξ is low.

• From P2's perspective, he faces a distribution of P1's actions.

Background: Mixed Strategies & Harsanyi Purification

Harsanyi: What if players' observe some private payoff perturbations?

-	L	R
Т	2, 1	0,0
В	0,0	1,2

Suppose P1's payoff from T equals his payoff in the matrix plus ξ .

• ξ is drawn from a continuous distribution $U[0, \varepsilon]$.

Only P1 observes the realization of ξ .

Suppose P2's payoff from L equals his payoff in the matrix plus η .

η is drawn from a continuous distribution U[0, ε].
 Only P2 observes the realization of η.

P2 chooses L when η is high and chooses R when η is low.

• From P1's perspective, he faces a distribution of P2's actions.

Results

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Purification in Dynamic Games

Background: Mixed Strategies & Harsanyi Purification

Let's solve for the Bayes Nash equilibrium in the perturbed game:

-	L	R
Т	$2+\xi, 1+\eta$	ξ ,0
В	0, η	1,2

Let $\xi^*, \eta^* \in [0, \varepsilon]$ be such that:

- P1 chooses *B* iff $\xi < \xi^* \Rightarrow$ P1 chooses *B* with prob $\frac{\xi^*}{\varepsilon}$.
- P2 chooses *R* iff $\eta < \eta^* \Rightarrow$ P2 chooses *R* with prob $\frac{\eta^*}{\varepsilon}$.

Type ξ^* player 1's indifference between *T* and *B* requires that

$$\Bigl(1-\frac{\eta^*}{\varepsilon}\Bigr)(2+\xi^*)=\frac{\eta^*}{\varepsilon}$$

$$\Rightarrow \quad \varepsilon(2+\xi^*) = \eta^*(3+\xi^*).$$

Results

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Purification in Dynamic Games

Background: Mixed Strategies & Harsanyi Purification

Let's solve for the equilibrium in the perturbed game:

-	L	R
Т	$2+\xi, 1+\eta$	ξ ,0
В	0, η	1,2

Let $\xi^*, \eta^* \in [0, \varepsilon]$ be such that:

- P1 chooses *B* iff $\xi < \xi^* \Rightarrow$ P1 chooses *B* with prob $\frac{\xi^*}{\varepsilon}$.
- P2 chooses *R* iff $\eta < \eta^* \Rightarrow$ P2 chooses *R* with prob $\frac{\eta^*}{\varepsilon}$.

Type η^* player 2's indifference between *L* and *R* requires that:

$$\Bigl(1-\frac{\xi^*}{\varepsilon}\Bigr)(1+\eta^*)=2\frac{\xi^*}{\varepsilon}$$

$$\Rightarrow \quad \varepsilon(1+\eta^*) = \xi^*(3+\eta^*).$$

Results

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Background: Mixed Strategies & Harsanyi Purification

Let's solve for the equilibrium in the perturbed game:

_	L	R
Т	$2+\xi, 1+\eta$	ξ ,0
В	0, <mark>η</mark>	1, 2

These two equations lead to a unique positive solution (ξ^*, η^*) :

 $\varepsilon(1+\eta^*) = \xi^*(3+\eta^*) \quad \varepsilon(2+\xi^*) = \eta^*(3+\xi^*).$

As $\varepsilon \to 0$, $\xi^* / \varepsilon \to 1/3$ and $\eta^* / \varepsilon \to 2/3$.

 As ε → 0, players' action distribution in the perturbed game converge to the mixed-strategy equilibrium in the unperturbed game. Purification in Normal Form Games

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Harsanyi Purification Theorem

Results

- *N* player finite normal-form game $\mathcal{G} \equiv (I, A, u)$, with $A \equiv \prod_{i=1}^{n} A_i$.
- Perturbation for player *i*: η_i ∈ ℝ<sup>|A_i|</sub> drawn according to continuous distribution μ_i s.t. player *i*'s payoff from *a* is u_i(a) + η_i(a).
 </sup>
- Only player *i* observes η_i , η_i is independent of η_j for every $i \neq j$.

Purification Theorem

Liu and Skrzypacz

For every regular Nash equilibrium $\sigma \in \Delta(A)$ of \mathcal{G} and for every

$$\{\mu_1^k, ..., \mu_n^k\}_{k \in \mathbb{N}}$$
 with $\lim_{k \to +\infty} \mu_i^k = 0$ for every *i*.

For every $\varepsilon > 0$ there exists $\overline{k} \in \mathbb{N}$ such that for every $k > \overline{k}$, the perturbed game $\{\mathcal{G}, \mu_1^k, ..., \mu_n^k\}$ has a Bayes Nash equilibrium s.t.

- each player has a strict incentive almost surely,
- this BNE induces a distribution that is within ε of σ .

Harsanyi Purification Theorem

Purification Theorem

For every regular Nash equilibrium $\sigma \in \Delta(A)$ of \mathcal{G} and for every $\{\mu_1^k, ..., \mu_n^k\}_{k \in \mathbb{N}}$ with $\lim_{k \to +\infty} \mu_i^k = 0$ for every *i*. For every $\varepsilon > 0$ there exists $\overline{k} \in \mathbb{N}$ such that for every $k > \overline{k}$, the perturbed game $\{\mathcal{G}, \mu_1^k, ..., \mu_n^k\}$ has a Bayes Nash equilibrium s.t.

- each player has a strict incentive almost surely,
- this BNE induces a distribution that is within ε of σ .

Regular Nash equilibrium is a refinement of Nash equilibrium.

• In generic finite normal-form games, every Nash equilibrium is a regular Nash equilibrium (Van Damme 1991).

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Definition of Regular Equilibrium

Fix a strategy profile
$$a^* \equiv (a_1^*, ..., a_n^*)$$
. For every $p \in \Delta(A)$, let

$$F_i^k(\boldsymbol{p}|a^*) \equiv p_i^k \Big\{ u_i(\boldsymbol{p}_{-i}, a_i^k) - u_i(\boldsymbol{p}_{-i}, a_i^*) \Big\} \text{ for every } i \in I \text{ and } a_i^k \neq a_i^*,$$

$$F_i(\mathbf{p}|a^*) = \sum_{k=1}^{|A_i|} p_i^k - 1$$

Let

$$J(\boldsymbol{p}^*, a^*) \equiv \frac{\partial \mathbf{F}(\boldsymbol{p}|a^*)}{\partial \boldsymbol{p}}\Big|_{\boldsymbol{p}=\boldsymbol{p}^*}$$

An equilibrium p^* is *regular* if and only if there exists $a^* \in A$ such that $J(p^*, a^*)$ is not a singular matrix.

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From Normal Form to Extensive Form

Can we generalize this insight to extensive-form games?

Bhaskar: Overlapping Generation Repeated Games

- Time t = 0, 1, 2, ... One agent born in each period.
- The agent who is born is period *t*:
 - Receives $K \in \mathbb{N}$ units of endowment in period *t*.
 - He shares $a_t \in \{0, 1, ..., K\}$ with his predecessor.
 - He may receive transfers from his successor in period t + 1.
- Agent *t*'s payoff is $u(a_t, a_{t+1})$, which is strictly increasing in a_{t+1} and is strictly decreasing in a_t .
- The efficient level of transfer:

$$k^* \in \arg \max_{k \in \{0,1,\dots,K\}} \Big\{ u(K-k,k) + u(k,K-k) \Big\}.$$

- We assume that k^* is strictly positive.
- We also assume that $u(K k^*, k^*) > u(K, 0)$.

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Benchmark: Unbounded Memories

Benchmark: Sustaining Cooperation with Perfect Information

If every agent can perfectly observe all previous agents' actions, then there exists a pure-strategy equilibrium that sustains transfer level k^* .

Proof: Grim-trigger strategies.

Introducing Limited Memories

For every j > i, agent j's info about agent i's action is a partition $B_{j,i}$ of A_i .

- Assumption 1: For every k > j > i, $B_{k,i}$ is weakly coarser than $B_{j,i}$.
- Assumption 2: There are infinitely many agents whose actions are no longer visible after a certain period.

Can we sustain cooperation via pure-strategy equilibria?

Theorem

Liu and Skrzypacz

Suppose $u(a_t, a_{t+1}) \neq u(a'_t, a'_{t+1})$ for any $(a_t, a_{t+1}) \neq (a'_t, a'_{t+1})$.

All agents choose 0 in every pure-strategy equilibrium.

Liu and	Skrzypacz
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Proof

Theorem

Suppose $u(a_t, a_{t+1}) \neq u(a'_t, a'_{t+1})$ for any $(a_t, a_{t+1}) \neq (a'_t, a'_{t+1})$. All agents choose 0 in every pure-strategy equilibrium.

Suppose agent *i*'s action is not visible starting from period j + 1.

- Will agent *j*'s action depend on agent *i*'s action?
- Suppose h_j and h'_j differ only in agent *i*'s action, and agent *j* plays *a* at h_j and plays *a'* at h'_j . Then agent *j*'s incentive constraints imply that:

$$u(a, a_{j+1}(h_j, a)) \ge u(a', a_{j+1}(h_j, a'))$$

and

$$u(a', a_{j+1}(h'_j, a')) \ge u(a, a_{j+1}(h'_j, a))$$

Not true since $a_{j+1}(h_j, a) = a_{j+1}(h'_j, a)$ and $a_{j+1}(h_j, a') = a_{j+1}(h'_j, a')$.

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Proof

Theorem

Suppose $u(a_t, a_{t+1}) \neq u(a'_t, a'_{t+1})$ for any $(a_t, a_{t+1}) \neq (a'_t, a'_{t+1})$. All players choose 0 in every pure-strategy equilibrium.

Suppose agent *i*'s action is not visible starting from period j + 1.

- Agent *j*'s action does not depend on agent *i*'s action.
- Can agent j 1's action depend on agent *i*'s action?
- Can agent *i* + 1's action depend on agent *i*'s action?

What will agent *i* choose when no agent's action depends on his? What will agent i - 1 choose when agent *i* chooses 0 no matter what?

Sustaining Cooperation via Mixed Strategies

Theorem

Suppose $u(K - k^*, k^*) > u(K, 0)$ and agent t + 1 observes the action of agent t. There exists a mixed strategy equilibrium that sustains transfer level k^* .

Consider an equilibrium in which

- Agent 0 chooses k^* .
- Agent *i* chooses k^* if agent i 1 chooses k^* , and mixes between 0 and k^* if agent i 1 chooses any other action.

Agent *i* is indifferent between 0 and k^* ,

- Their mixing probability depends on agent i 1's action, and is chosen in order to make agent i 1 indifferent.
- Agent *i*'s incentive to mix is provided by agent i + 1's mixed action.

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Are These Mixed-Strategy Equilibria Robust?

Suppose agents' payoffs are perturbed so that agent *t*'s payoff from (a_t, a_{t+1}) equals:

$$u(a_t, a_{t+1}) + \eta_t(a_t, a_{t+1}),$$

where $\eta_t \equiv {\eta_t(a_t, a_{t+1})}_{(a_t, a_{t+1}) \in A^2}$ is drawn according to continuous distribution F^n , and the support of $\eta_t(a_t, a_{t+1})$ is close to 0.

- η_t is independent of η_s for every $t \neq s$,
- η_t is independent of the history of play.

Theorem

When the support of the perturbations are close to 0, all agents choose 0 in every equilibrium of every perturbed game.

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Proof Sketch

Theorem

All agents choose 0 in every equilibrium of every perturbed game.

Suppose agent *i*'s action is not visible starting from period j + 1.

- Will agent *j*'s action depend on agent *i*'s action?
- Agent *j* strictly prefers *a* to a' at h_j iff:

 $u(a, a_{j+1}(h_j, a)) + \eta_j(a, a_{j+1}(h_j, a)) > u(a', a_{j+1}(h_j, a')) + \eta_j(a', a_{j+1}(h_j, a')).$

or equivalently,

 $\eta_j(a, a_{j+1}(h_j, a)) - \eta_j(a', a_{j+1}(h_j, a')) > u(a', a_{j+1}(h_j, a')) - u(a, a_{j+1}(h_j, a)).$

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Proof Sketch

Theorem

All agents choose 0 in every equilibrium of every perturbed game.

Suppose agent *i*'s action is not visible starting from period j + 1.

- Suppose h'_i and h_j differ only in agent *i*'s action
- Agent *j* strictly prefers *a* to a' at h'_j iff:

 $u(a, a_{j+1}(h'_j, a)) + \eta_j(a, a_{j+1}(h'_j, a)) > u(a', a_{j+1}(h'_j, a')) + \eta_j(a', a_{j+1}(h'_j, a')).$ or equivalently,

 $\eta_j(a, a_{j+1}(h_j', a)) - \eta_j(a', a_{j+1}(h_j', a')) > u(a', a_{j+1}(h_j', a')) - u(a, a_{j+1}(h_j', a)).$

Proof Sketch

Theorem

All agents choose 0 in every equilibrium of every perturbed game.

Suppose agent *i*'s action is not visible starting from period j + 1.

• Suppose h'_j and h_j differ only in agent *i*'s action.

Then $a_{j+1}(h_j, a) = a_{j+1}(h'_j, a)$ and $a_{j+1}(h_j, a') = a_{j+1}(h'_j, a')$.

• Agent *j* strictly prefers *a* to *a'* at *h_j* iff:

 $\eta_j(a, a_{j+1}(h_j, a)) - \eta_j(a', a_{j+1}(h_j, a')) > u(a', a_{j+1}(h_j, a')) - u(a, a_{j+1}(h_j, a)).$

• Agent *j* strictly prefers *a* to a' at h'_j iff:

 $\eta_j(a, a_{j+1}(h_j', a)) - \eta_j(a', a_{j+1}(h_j', a')) > u(a', a_{j+1}(h_j', a')) - u(a, a_{j+1}(h_j', a)).$

• Prob that agent *j* chooses each action is independent of agent *i*'s action.