Lecture 14: Behavioral Social Learning

Harry PEI Department of Economics, Northwestern University

Spring Quarter, 2021

Motivation

Social learning with homogenous preferences and rational Bayesian agents.

 Agents' actions are asymptotically efficient if their private signals are unbounded (Smith and Sorensen), or their action space is sufficiently rich (Lee).

Critiques of rational social learning models.

• Requires too much sophistication (e.g., double-counting problems).

Attempts to relax the rationality assumption.

- DeGroot (1974), Golub and Jackson (2010), Molavi et al. (2012,2018): Non-Bayesian rule-of-thumb learning rules.
- Eyster and Rabin: Agents are Bayesian but fails to recognize the double-counting problem when aggregating different sources of info.

Social Learning in a Doubly Rich Setting

Social learning with unbounded signals and a continuum of actions.

- Time t = 1, 2, ... One agent arriving in each period.
- State $\omega \in \{0,1\}$, equally likely.
- Action of agent $t: a_t \in [0, 1]$.
- Agent *t* observes s_t and $\{a_1, ..., a_{t-1}\}$ and then chooses a_t .
- Agent *t*'s payoff is $-(a_t \omega)^2$, so $a_t \equiv \mathbb{E}[\omega|s_t, a_1, ..., a_{t-1}]$.
- Agent *t*'s private signal $s_t \sim G(\cdot | \omega)$, conditionally independent.
 - ↔ The signal structure can be represented by the distribution over private beliefs conditional on ω, i.e., $F_ω ∈ Δ[0,1]$.
 - \hookrightarrow We assume that for every ω , $F_{\omega}(0) = 0$, $F_{\omega}(1) = 1$, F_{ω} is differentiable, and has continuous and positive density f_{ω} .
 - \hookrightarrow Unbounded private beliefs, no perfectly revealing signal.
- Let $p_t \equiv \mathbb{E}[\boldsymbol{\omega}|s_t]$.

The log likelihood ratio (LLR) is $l_t \equiv \log \frac{p_t}{1-p_t}$.

Benchmark: Sophisticated Bayesian Social Learning

Suppose agents are Bayesian and rational (i.e., fully sophisticated).

Theorem: Smith and Sorensen (2000)

Suppose the agents' private signals are unbounded, then conditional on every $\omega \in \{0,1\}$, $a_t \rightarrow \omega$ almost surely.

Theorem: Lee (1993)

In environments where $a_t \in [0,1]$, then conditional on every $\omega \in \{0,1\}$, $a_t \rightarrow \omega$ almost surely.

Sophisticated Bayesian agents' actions are asymptotically efficient.

Naive Bayesian Agents

Form of naivete: Best response trailing naive inference (BRTNI)

- Each player best responds to the belief that each of her predecessors follows their own signal.
- This reasoning neglects the fact that their predecessors also make inferences from their own predecessors' actions.

What are the beliefs of these naive players?

- Player 1's posterior log likelihood ratio is l_1 .
- Player 2's posterior log likelihood ratio is $l_1 + l_2$.
- Player 3's posterior log likelihood ratio is $2l_1 + l_2 + l_3$.

If P3 is rational, then they would ignore P1's action and their posterior should be $l_1 + l_2 + l_3$.

• Player 4's posterior log likelihood ratio is $4l_1 + 2l_2 + l_3 + l_4$.

Naive Bayesian Agents

... ...

What are players' beliefs if they engage in BRTNI?

- Player 1's posterior log likelihood ratio is l_1 .
- Player 2's posterior log likelihood ratio is $l_1 + l_2$.
- Player 3's posterior log likelihood ratio is $2l_1 + l_2 + l_3$.
- Player 4's posterior log likelihood ratio is $4l_1 + 2l_2 + l_3 + l_4$.
- Player *n*'s posterior log likelihood ratio is $l_n + \sum_{\tau=1}^{n-1} 2^{n-1-\tau} l_{\tau}$.

Players over-weight the private signals of early players.

- P1's private signal s_1 should have weight 1/t in Player t's belief.
- When players are naive, s_1 has weight 1/2 in all players' beliefs.

It is hard to correct early players' mistakes even with rich action spaces and unbounded private signals (will affect the asymptotic outcome).

Inefficiencies in All Periods

Theorem

When players are naive, for every r < 1, there exists $\delta > 0$ such that

$$\Pr\left(a_t > r \text{ for all } t \in \mathbb{N} \middle| \boldsymbol{\omega} = 0\right) > \boldsymbol{\delta}.$$

Naive players' actions are bounded away from efficiency in *all periods* with positive probability.

• Intuition: Since naive agents over-weigh the signals of earlier agents, it is hard to correct earlier players' mistakes.

Next: How the proof incorporates this intuition.

Proof

Theorem

When players are naive, for every r < 1, there exists $\delta > 0$ such that

$$\Pr\left(a_t > r \text{ for all } t \in \mathbb{N} \middle| \boldsymbol{\omega} = 0\right) > \boldsymbol{\delta}.$$

Let L_t be the log likelihood ratio after a naive player observes all predecessors' actions but before observing their own private signal.

• Given that the prior is uniform, $\log \frac{a_t}{1-a_t} = L_t + l_t$.

One can show by induction that $L_n = 2L_{n-1} + l_{n-1}$ for every $n \in \mathbb{N}$.

- Public LLR in period $n-1: \sum_{\tau=1}^{n-2} \log \frac{a_{\tau}}{1-a_{\tau}}$.
- Player n-1's action satisfies $\log \frac{a_{n-1}}{1-a_{n-1}} = L_{n-1} + l_{n-1}$.
- Public LLR in period *n* is $L_{n-1} + \log \frac{a_{n-1}}{1-a_{n-1}} = 2L_{n-1} + l_{n-1}$.

Proof

Theorem

When players are naive, for every r < 1, there exists $\delta > 0$ such that

$$\Pr\left(a_t > r \text{ for all } t \in \mathbb{N} \middle| \boldsymbol{\omega} = 0\right) > \boldsymbol{\delta}.$$

When players are naive, L_n satisfies $L_n = 2L_{n-1} + l_{n-1}$.

- Pick an arbitrary $r \in (1/2, 1)$, let $R \equiv \log \frac{r}{1-r} > 0$.
- Question: What if $L_2 > 3R$ and $l_t \ge -tR$ for every $t \in \mathbb{N}$?

Since $L_2 = l_1$, LLR of P1's action is greater than 3*R*.

Since P2's LLR is $L_2 + l_2$ and $l_t \ge -2R$, we have $\log \frac{a_2}{1-a_2} > R$.

Since $L_3 = 2L_2 + l_2 \ge 6R - 2R = 4R$ and $l_3 \ge -3R$, $\log \frac{a_3}{1 - a_2} > R$.

Lemma

If $L_2 > 3R$ and $l_t \ge -tR$ for every $t \in \mathbb{N}$, then $L_n > (n+1)R$ and $\log \frac{a_n}{1-a_n} > R$ for every $n \ge 2$ (which implies that $a_t > r$).

Proof

Theorem

When players are naive, for every r < 1, there exists $\delta > 0$ such that

$$\Pr\left(a_t > r \text{ for all } t \in \mathbb{N} \middle| \boldsymbol{\omega} = 0\right) > \boldsymbol{\delta}.$$

Lemma

If $L_2 > 3R$ and $l_t \ge -tR$ for every $t \in \mathbb{N}$, then $L_n > (n+1)R$ and $\log \frac{a_n}{1-a_n} > R$ for every $n \ge 2$.

Intuition behind the lemma: Suppose $\omega = 0$,

- $L_2 \ge 3R$ means that P1's signal is in favor of $\omega = 1$.
- *l_t* ≥ −*tR*: mistake of P1's signal becomes harder to correct over time. Why? Since P1's signal carries a large weight, *l_t* needs to be sufficiently negative in order to drive log ^{*a_t*}/_{1−a_t} below *R*.

Proof

Theorem

When players are naive, for every r < 1, there exists $\delta > 0$ such that

$$\Pr\left(a_t > r \text{ for all } t \in \mathbb{N} \,\middle|\, \boldsymbol{\omega} = 0\right) > \boldsymbol{\delta}.$$

Lemma

If $L_2 > 3R$ and $l_t \ge -tR$ for every $t \in \mathbb{N}$, then $L_n > (n+1)R$ and $\log \frac{a_n}{1-a_n} > R$ for every $n \ge 2$.

Since $L_2 > 3R$ with positive prob conditional on $\omega = 0$, we only need to show that event

$$\left\{l_t \geq -tR \text{ for every } t \in \mathbb{N}\right\}$$

occurs with prob bounded away from 0 conditional on $\omega = 0$.

Proof

We need to show that

$$\Pr\left\{l_t \geq -tR \text{ for every } t \in \mathbb{N} \middle| \boldsymbol{\omega} = 0\right\} > 0.$$

By Markov inequality,

$$\Pr\left\{l_t < -tR \left| \boldsymbol{\omega} = 0\right\} \le \Pr\left\{l_t^2 \ge (tR)^2 \left| \boldsymbol{\omega} = 0\right\} \le \frac{1}{(tR)^2} \mathbb{E}[l_t^2 | \boldsymbol{\omega} = 0].$$

Bound the value of $Q \equiv \mathbb{E}[l_t^2 | \omega = 0]$ from above:

$$\mathbb{E}[l_t^2|\boldsymbol{\omega}=0] = \int_0^1 \left(\log\frac{s}{1-s}\right)^2 f_0(s) ds \le \underbrace{\max\{f_0(s)|s\in[0,1]\}}_{\text{a bounded number}} \underbrace{\int_0^1 \left(\log\frac{s}{1-s}\right)^2 ds}_{=\pi^2/3}$$

Hence, there exists a bounded Q such that for every $t \in \mathbb{N}$,

$$\Pr\left\{l_t \ge -tR \middle| \omega = 0\right\} \ge 1 - \frac{Q}{(tR)^2}.$$

Proof

We need to show that

$$\Pr\left\{l_t \geq -tR \text{ for every } t \in \mathbb{N} \middle| \boldsymbol{\omega} = 0\right\} > 0.$$

We have shown that there exists Q > 0 such that for every $t \in \mathbb{N}$,

$$\Pr\left\{l_t \geq -tR \middle| \boldsymbol{\omega} = 0\right\} \geq 1 - \frac{Q}{(tR)^2}.$$

Let $\tau \in \mathbb{N}$ be such that $1 - \frac{Q}{(tR)^2} > 0$ for every $t > \tau$.

- We know that $\Pr\left\{l_t \ge -tR \text{ for every } t \le \tau \middle| \omega = 0\right\} > 0.$
- We need to show that $\Pr\left\{l_t \ge -tR \text{ for every } t > \tau \mid \omega = 0\right\} > 0.$

Proof

We have shown that for every $t \in \mathbb{N}$,

$$\Pr\left\{l_t \geq -tR \middle| \, \omega = 0\right\} \geq 1 - \frac{Q}{(tR)^2}.$$

We need to show that $\Pr\left\{l_t \ge -tR \text{ for every } t > \tau \middle| \omega = 0\right\} > 0.$

$$\Pr\left\{l_t \ge -tR \text{ for every } t > \tau \middle| \omega = 0\right\} \ge \Pi_{t=\tau+1}^{+\infty} \left(1 - \frac{Q}{(tR)^2}\right)$$
$$= \underbrace{\exp\left\{\sum_{t>\tau} \log\left(1 - \frac{Q}{(tR)^2}\right)\right\} \ge \exp\left\{\sum_{t>\tau} - \frac{Q}{(tR)^2}\right\}}_{\text{uses inequality } \log(1-x)\ge -x} \ge \exp\left(-\frac{Q\pi}{6R^2}\right)$$

To summarize, both $\Pr \left\{ l_t \ge -tR \text{ for every } t > \tau \middle| \omega = 0 \right\}$ and $\Pr \left\{ l_t \ge -tR \text{ for every } t \le \tau \middle| \omega = 0 \right\}$ are bounded away from 0, which implies that $\Pr \left\{ l_t \ge -tR \text{ for every } t \middle| \omega = 0 \right\}$ is bounded away from 0.

Limit Points of Naive Agents' Actions

Theorem

When agents are naive, for every r < 1, there exists $\delta > 0$ such that

$$\Pr\left(a_t > r \text{ for all } t \in \mathbb{N} \middle| \boldsymbol{\omega} = 0\right) > \boldsymbol{\delta}.$$

The agents' limiting actions can be wrong, but what can they be?

Theorem

When agents are naive, their beliefs (and hence their actions) converge almost surely to either 0 or 1.

Lesson: If their beliefs are wrong in the long run, then they must be fully confident in the wrong state.

• Cannot happen when agents are Bayesian and sophisticated.

Proof

To show that players' beliefs converge a.s. to 0 or 1, it is sufficient to show that L_n diverges to $+\infty$ or $-\infty$ almost surely as $n \to +\infty$.

Recall the formula for L_n :

$$L_n = \sum_{t=1}^{n-1} \log \frac{a_t}{1-a_t} = \sum_{t=1}^{n-1} 2^{n-t-1} l_t$$

Therefore,

$$2^{1-n}L_n = \sum_{t=1}^{n-1} 2^{-t} l_t.$$

If we can show that $\sum_{t=1}^{n-1} 2^{-t} l_t$ converges as $n \to +\infty$, then L_n must be diverging to $+\infty$ or $-\infty$.

Proof

We need to show that $\sum_{t=1}^{n-1} 2^{-t} l_t$ converges as $n \to +\infty$.

Kolmogorov Three-Series Theorem (Theorem 5.3.3 in Chung's textbook)

Suppose $\{X_n\}_{n\in\mathbb{N}}$ are independent random variables. Then $\sum_n X_n$ converges *a.s.* if the following conditions hold for some A > 0

- 1. $\sum_{n} \Pr(|X_n| \ge A)$ converges,
- 2. $\sum_{n} \mathbb{E}[X_n \mathbf{1}\{|X_n| \le A\}]$ converges,
- 3. $\sum_{n} Var(X_n \mathbf{1}\{|X_n| \le A\})$ converges.

Let X_n be $2^{-n}l_n$ conditional on $\omega = 0$.

$$\sum \Pr(|X_n| \ge A | \boldsymbol{\omega} = 0) = \sum \Pr(2^{-n} | l_n | \ge A | \boldsymbol{\omega} = 0) = \sum \Pr(4^{-n} l_n^2 \ge A^2 | \boldsymbol{\omega} = 0)$$

$$\leq \sum \frac{4^{-n} \mathbb{E}[l_n^2 | \boldsymbol{\omega} = 0]}{A^2} \leq \frac{\mathbb{E}[l_n^2 | \boldsymbol{\omega} = 0]}{A^2} \quad \text{(which is bounded)}$$

Proof

Kolmogorov Three-Series Theorem (Theorem 5.3.3 in Chung's textbook)

Suppose $\{X_n\}_{n\in\mathbb{N}}$ are independent random variables. Then $\sum_n X_n$ converges *a.s.* if the following conditions hold for some A > 0

- 1. $\sum_{n} \Pr(|X_n| \ge A)$ converges,
- 2. $\sum_{n} \mathbb{E}[X_n \mathbf{1}\{|X_n| \le A\}]$ converges,
- 3. $\sum_{n} Var\left(X_n \mathbf{1}\{|X_n| \le A\}\right)$ converges.

Let X_n be $2^{-n}l_n$ conditional on $\omega = 0$.

$$\sum \mathbb{E}\left[2^{-n}l_n \mathbf{1}\left\{2^{-n}l_n \le A\right\} \middle| \boldsymbol{\omega} = 0\right] \le \sum \mathbb{E}\left[2^{-n}|l_n| \middle| \boldsymbol{\omega} = 0\right] \le \sum \mathbb{E}\left[2^{-n}(l_n^2 + 1) \middle| \boldsymbol{\omega} = 0\right]$$

 $=\sum 2^{-n}+\sum 2^{-n}\mathbb{E}[l_n^2|\boldsymbol{\omega}=0]=1+\mathbb{E}[l_n^2|\boldsymbol{\omega}=0] \quad (\text{which is bounded}).$

Proof

Kolmogorov Three-Series Theorem (Theorem 5.3.3 in Chung's textbook)

Suppose $\{X_n\}_{n\in\mathbb{N}}$ are independent random variables. Then $\sum_n X_n$ converges *a.s.* if the following conditions hold for some A > 0

- 1. $\sum_{n} \Pr(|X_n| \ge A)$ converges,
- 2. $\sum_{n} \mathbb{E}[X_n \mathbf{1}\{|X_n| \le A\}]$ converges,
- 3. $\sum_{n} Var\left(X_n \mathbf{1}\{|X_n| \le A\}\right)$ converges.

Let X_n be $2^{-n}l_n$ conditional on $\omega = 0$.

$$\sum \operatorname{Var}\left(X_n \mathbf{1}\{|X_n| \le A\} \middle| \boldsymbol{\omega} = 0\right) \le \sum \mathbb{E}\left[4^{-n} l_n^2 \middle| \boldsymbol{\omega} = 0\right] \le \mathbb{E}[l_n^2 \middle| \boldsymbol{\omega} = 0].$$

The convergence of all three series uses the fact that $\mathbb{E}[l_n^2|\omega=0]$ is bounded.

• Hinges on the existence of continuous density f_{ω} .

Stable Interior Beliefs are Likely to Be Wrong

Suppose players' beliefs remain stable at some interior level for a long time, what happens?

Theorem

For every $[c,d] \subset (1/2,1)$, there exists $T \in \mathbb{N}$ such that if $a_t \in [c,d]$ for every $t \in \{1,2,...,T\}$, then

$$\Pr\left(\omega=0\Big|(a_1,...,a_T)\right)>\Pr\left(\omega=1\Big|(a_1,...,a_T)\right)$$

Why is $\omega = 0$ more likely to be the correct state when agents' belief stablize at an interval above 1/2?

- Suppose $a_1 \in [c,d] \subset (1/2,1)$. If $\omega = 1$, then $l_2, l_3, ..., l_n$ are likely to be high, which means that a_i will approach 1 in the long run.
- Hence, $a_t \in [c,d]$ for a long time indicates that ω is likely to be 0.

Stable Interior Beliefs are Likely to Be Wrong

Theorem

For every $[c,d] \subset (1/2,1)$, there exists $T \in \mathbb{N}$ such that if $a_t \in [c,d]$ for every $t \in \{1,2,...,T\}$, then

$$\Pr\left(\boldsymbol{\omega}=0\middle|(a_1,...,a_T)\right)>\Pr\left(\boldsymbol{\omega}=1\middle|(a_1,...,a_T)\right)$$

Let
$$u \equiv \log \frac{c}{1-c}$$
 and $v \equiv \log \frac{d}{1-d}$.

• If $\log \frac{a_1}{1-a_1}, \dots, \log \frac{a_t}{1-a_t} \in [u, v]$, then

$$v \ge \log \frac{a_{t+1}}{1 - a_{t+1}} = \sum_{\tau=1}^{t} \log \frac{a_{\tau}}{1 - a_{\tau}} + l_{t+1} \ge tu + l_{t+1}$$

which means that

$$l_{t+1} \leq v - tu.$$

- When *t* is large enough, s_{t+1} is a signal in favor of state 0.
- When $l_{t+1} \le v tu$ for all $t \le T$ and T being large enough, the posterior prob of $\omega = 0$ under a rational agent's belief is less than 1/2.

Network Among Players

Doubly rich setting, agents are Bayesian, but believe that all their predecessors' actions only reflect their own private signals.

- Agents' belief converges to the wrong state with positive prob.
- Agents' actions are asymptotically inefficient with positive prob.

How general is this finding?

- What if players cannot observe all their predecessors' actions?
- What if players exhibit redundancy neglect, but not as extreme as in the previous model?

Eyster and Rabin (2014)

- Time t = 1, 2, ... One agent arriving in each period.
- State $\omega \in \{0, 1\}$, equally likely.
- Action of agent $t: a_t \in [0, 1]$.
- Agent *t*'s payoff is $-(a_t \omega)^2$, so $a_t \equiv \mathbb{E}[\omega|s_t, a_1, ..., a_{t-1}]$.
- Agent *t*'s private signal $s_t \sim G(\cdot | \omega)$, conditionally independent.

 \hookrightarrow same assumption as before (unbounded, no revealing signal).

- Agent *t* observes s_t and $\{a_{\tau}\}_{\tau \in N_t}$ where $N_t \subset \{1, 2, ..., t-1\}$.
 - \hookrightarrow Players in N_t are the neighbors of agent t.
 - $\,\hookrightarrow\,$ The network is deterministic and is common knowledge.
- We write $i \succ j$ if there exist k(0), ..., k(n) such that k(0) = j, k(n) = i, $k(m-1) \in N_{k(m)}$ for every $m \in \{1, 2, ..., n\}$.

Player *i* can observe player *j*'s action along some path.

Strategies & Regularity Assumptions on Strategies

Let l_t be the LLR of agent *t*'s private signal and let $\alpha_t \equiv \log \frac{a_t}{1-a_t}$.

Agent *t*'s strategy is
$$\alpha_t(\alpha_1, ..., \alpha_{t-1}, l_t)$$
, measurable w.r.t $(l_t, (\alpha_\tau)_{\tau \in N(t)})$.

Strictly Increasing Strategies

Players' strategies are strictly increasing in private signals if for every $t \in \mathbb{N}$ *and* $(\alpha_1, ..., \alpha_{t-1})$ *,* α_t *is a strictly increasing function of* l_t *.*

Boundedly Increasing Strategies

Players' strategies are boundedly increasing if there exists $K \in \mathbb{R}_+$ *such that for every* $t \in \mathbb{N}$ *,* $(\alpha_1, ..., \alpha_{t-1})$ *, and* $l_t \neq l'_t$ *, we have*

$$\left| lpha_t \Big(lpha_1, ..., lpha_{t-1}, l_t \Big) - lpha_t \Big(lpha_1, ..., lpha_{t-1}, l_t' \Big) \right| \leq K |l_t - l_t'|$$

General Redundancy Neglect Learning Rules

Redundancy Neglect Strategies

Players' strategies exhibit redundancy neglect if

- 1. For every t and $j \prec t$, α_t is weakly increasing in α_j regardless of $(\alpha_1, ..., \alpha_{t-1})$ and l_t .
- 2. There exist $N \in \mathbb{N}$ and x > 1 such that for every player $t \ge N+1$ and z' > z, α_t increases by at least x(z'-z) if each of $\alpha_{t-N}, ..., \alpha_{t-1}$ increases from z to z'.

This is a joint condition on the network and players' learning rule.

• Player t observes at least one of their last N predecessors.

The learning rule in the previous model satisfies both requirements.

•
$$\alpha_t = l_t + \sum_{\tau=1}^{t-1} \alpha_{\tau}$$
.

How does rational Bayesian learning violate these two properties?

General Redundancy Neglect Learning Rules

Redundancy Neglect Strategies

Players' strategies exhibit redundancy neglect if

- 1. For every t and $j \prec t$, α_t is weakly increasing in α_j regardless of $(\alpha_1, ..., \alpha_{t-1})$ and l_t .
- 2. There exists $N \in \mathbb{N}$ and x > 1 such that for every player $t \ge N+1$ and z' > z, α_t increases by at least x(z'-z) if each of $\alpha_{t-N}, ..., \alpha_{t-1}$ increases from z to z'.

Suppose there are four individuals.

- 2 and 3 can both observe 1, but cannot observe each other.
- 4 observes 1, 2, and 3.

1's action is l_1 , 2's action is $\alpha_1 + l_2$, 3's action is $\alpha_1 + l_3$.

- 4's optimal action is $l_1 + l_2 + l_3 + l_4$.
- His strategy is $l_4 + \alpha_2 + \alpha_3 \alpha_1$ (violates requirement 1).

General Redundancy Neglect Learning Rules

Redundancy Neglect Strategies

Players' strategies exhibit redundancy neglect if

- 1. For every t and $j \prec t$, α_t is weakly increasing in α_j regardless of $(\alpha_1, ..., \alpha_{t-1})$ and l_t .
- 2. There exists $N \in \mathbb{N}$ and x > 1 such that for every player $t \ge N + 1$, if each of $\alpha_{t-N}, ..., \alpha_{t-1}$ increases by at least Δ , then α_t increases by at least $x\Delta$.

Suppose every agent can observe all their predecessors.

• Agent *n*'s optimal action $\alpha_n = \alpha_{n-1} + l_n$ (violates requirement 2).

Theorem

If players' strategies are strictly and boundedly increasing, and exhibit redundancy neglect, then

- conditional on $\omega = 0$, α_t converges to $+\infty$ with positive prob,
- conditional on $\omega = 1$, α_t converges to $-\infty$ with positive prob.

The proof uses ideas similar to that of their earlier result.

• Since players over-react to earlier players' private signals, early players' mistakes are hard to correct,

so incorrect actions can be taken asymptotically with positive prob.

When will rational players anti-imitate?

Suppose that players are rational and Bayesian.

• Player *t* anti-imitates player *j* if $t \succ j$ and α_t is a strictly decreasing function of α_j .

Players i, j, k, l form a shield if

- *j* and *k* observe *i*,
- *j* and *k* cannot observe each other,
- *l* observes *i*, *j*, and *k*.

Theorem

Suppose all players are rational. There exists anti-imitation in equilibrium if and only if the network contains a shield.

Anti-imitation cannot arise under the canonical observation structure.